

EXAM I SOLUTIONS, MA4027, Summer 2004

Needless to say, there are other solutions.

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1. A *strongly regular graph*, or *SR-graph*, is a graph  $G$  that is regular of degree  $k$ , with  $v$  vertices, satisfying the following conditions:

If  $p, q$  are adjacent vertices, then exactly  $\lambda$  vertices are adjacent to both  $p$  and  $q$ .

If  $p, q$  are nonadjacent vertices, then exactly  $\mu$  vertices are adjacent to both  $p$  and  $q$ .

Such a graph  $G$  is said to have parameters  $(v, k, \lambda, \mu)$ . If both  $G$  and its complement are connected,  $G$  is said to be a nontrivial *SR-graph*, otherwise  $G$  is a trivial *SR-graph*. Show that the Petersen graph is nontrivial *SR*, and find its parameters.

**Solution:** The Petersen Graph has ten vertices and is regular of degree three, so  $v = 10$  and  $k = 3$ . The neighborhoods of adjacent vertices are disjoint, so  $\lambda = 0$ , while each nonadjacent pair has a common neighbor, so  $\mu = 1$ . Letting  $G = (V, E)$  denote the Petersen graph, we now show that  $\overline{G}$  is connected. Let  $x, y \in V$ , and assume that  $xy \notin \overline{E}$ . Since  $|N(x)| = |N(y)| = 6$  and  $|N(x) \cup N(y)| \leq 8$ , it follows that  $N(x) \cap N(y) \neq \emptyset$ . Thus the Petersen graph is nontrivial *SR* with parameters  $v = 10$ ,  $k = 3$ ,  $\lambda = 0$ , and  $\mu = 1$ .

2. A saturated hydrocarbon is a molecule  $C_mH_n$  in which every carbon atom has four bonds, every hydrogen atom has one bond, and no sequence of bonds is cyclic. Give a graph-theoretic argument that, for every integer  $m$ ,  $C_mH_n$  can exist only if  $n = 2m + 2$ .

**Proof:** We can model a hypothetical saturated hydrocarbon with a graph  $G$  on  $m + n$  vertices. Since carbon atoms have four bonds, we have  $m$  vertices of degree four, while the remaining  $n$  vertices represent hydrogen atoms and have degree one. The graph is a tree, since no sequence of bonds is cyclic. Since  $G$  is a tree, the number of edges in  $E(G)$  is  $m + n - 1$ . By the degree-sum formula, the number of edges in  $E(G)$  is  $\frac{1}{2}(4m + n) = 2m + n/2$ . Thus  $m + n - 1 = 2m + n/2$ , and it follows that  $n = 2m + 2$ .  $\square$

3. Recall that a tournament is an orientation of  $K_n$ . A tournament  $G$  is *transitive* if, for all  $x, y, z \in V(G)$ ,  $xy \in E(G)$  and  $yz \in E(G)$  imply  $xz \in E(G)$ ; in other words, the adjacency relation is transitive. Prove that  $G$  is a transitive tournament if and only if  $G$  contains no directed cycle.

**Proof:** Let  $G$  be a tournament. By definition,  $G$  is transitive if and only if  $G$  contains no cyclic triple. First suppose that  $G$  contains a cycle of length  $k > 3$ . If  $k = 3$ , we're done, so assume  $k > 3$ . Suppose  $C$  is a directed cycle on vertices  $v_1, v_2, \dots, v_k$ . Let  $j$  be the first index in  $3, 4, \dots, k$  such that  $v_jv_1 \in E(G)$ . (Note that  $j$  is well-defined, since

$v_k v_1 \in E(G)$ .) Then  $v_1 v_{j-1} v_j v_1$  is a directed 3-cycle. It follows that if  $G$  contains a directed cycle of any length, then  $G$  contains a cyclic triple and is not transitive. Now suppose that  $G$  contains no cycles of any length. Then clearly  $G$  contains no cyclic triple, and thus  $G$  is transitive.  $\square$

4. Prove that  $G$  is bipartite if and only if every subgraph  $H$  of  $G$  contains an independent set consisting of at least half of  $V(H)$ .

**Proof:** First suppose that  $G = (X, Y, E)$  is bipartite. Let  $H = (X', Y', E')$  be a subgraph of  $G$ . Then  $H$  is bipartite. The larger of  $X', Y'$  (choose arbitrarily if  $|X'| = |Y'|$ ) is an independent set containing at least half of  $V(H)$ . Now suppose that  $G$  is not bipartite. Then  $G$  contains a subgraph  $H \cong C_{2k+1}$ . Since any set of more than  $k$  vertices of  $H$  must contain two adjacent vertices, no independent set in  $H$  contains half of  $V(H)$ .  $\square$

5. Use the König-Egerváry Theorem (Theorem 3.1.16) to prove Hall's Theorem (Theorem 3.1.11).

**Solution:** The theorem to be proven states that, for any bipartite graph  $G = (X, Y, E)$ ,  $G$  has an  $X$ -saturating matching iff  $|N(S)| \geq |S|$  for every  $S \subseteq X$ . The König-Egerváry Theorem states that if  $G = (X, Y, E)$  is bipartite, then the size of a smallest vertex cover is equal to that of a largest matching.

We now proceed with the

**Proof:** Let  $G = (X, Y, E)$ . If  $|N(S)| < |S|$  for some  $S \subseteq X$ , then no  $S$ -saturating matching, and therefore no  $X$ -saturating matching, can exist. We must prove the converse. Suppose, then, that  $|N(S)| \geq |S|$  for every  $S \subseteq X$ . It suffices to show that  $|K| \geq |X|$  for every cover  $K$ , since then by the König-Egerváry Theorem  $|M| \geq |X|$  for every maximum-cardinality matching  $M$ . So let  $K$  be a cover in  $G$ , and let  $S = X - K$ . Since  $K$  is a cover, and since  $S \not\subseteq K$ , then we know that  $N(S) \subseteq K$ . It follows that

$$|K| \geq |X - S| + |N(S)| = |X| - |S| + |N(S)| \geq |X|,$$

and the proof is complete.  $\square$